

Parallel Matlab programming using Distributed Arrays

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Slide-1 Parallel MATLAB



- In the past, writing well performing parallel programs has required a lot of code and a lot of expertise
- pMatlab distributed arrays eliminates the coding burden
 - However, making programs run fast still requires expertise
- This talk illustrates the key math concepts experts use to make parallel programs perform well





Outline



- Distributed Arrays
- Concurrency vs Locality
- Execution
- Summary

- Serial Program
- Parallel Execution
- Distributed Arrays
- Explicitly Local



Serial Program

<u>Math</u>	<u>Matlab</u>
	X = zeros(N,N);
$\mathbf{X,Y}: \mathbb{R}^{N imes N}$	Y = zeros(N,N);
$\mathbf{V} = \mathbf{V} + 1$	
$\mathbf{Y} = \mathbf{X} + 1$	Y(:,:) = X + 1;
Matlab is a high lev	vel language
 Allows mathematical expressions to be written concisely 	
Multi-dimensional arrays are <i>fundamental</i> to Matlab	



Parallel Execution



- Run N_P (or Np) copies of same program
 - Single Program Multiple Data (SPMD)
- Each copy has a unique P_{ID} (or Pid)
- Every array is replicated on each copy of the program



Distributed Array Program



- Use P() notation (or map) to make a distributed array
- Tells program which dimension to distribute data
- Each program implicitly operates on only its own data (owner computes rule)



Explicitly Local Program

Math	pMatlab
	piliatian
	$XYmap = map([Np 1], {}, 0:Np-1);$
	<pre>Xloc = local(zeros(N,N,XYmap));</pre>
$\mathbf{X}, \mathbf{Y} : \mathbb{R}^{P(N) \times N}$	<pre>Yloc = local(zeros(N,N,XYmap));</pre>
,	
V 71 V 71 . 1	
$\mathbf{Y}.\mathbf{loc} = \mathbf{X}.\mathbf{loc} + 1$	Yloc(:,:) = Xloc + 1;
• Use loc notation (o	r local function) to explicitly retrieve
	, <u> </u>
local part of a distr	-
 Operation is the same as serial program, but with different 	
data on each processor (recommended approach)	



Outline

Parallel Design



- Concurrency vs Locality
- Execution
- Summary



Parallel Data Maps

<u>Array</u>	<u>Math</u>	<u>Matlab</u>
	$\mathbb{R}^{P(N) extbf{x} N}$	<pre>Xmap=map([Np 1],{},0:Np-1)</pre>
	$\mathbb{R}^{N x P(N)}$	<pre>Xmap=map([1 Np],{},0:Np-1)</pre>
	$\mathbb{R}^{P(N) x P(N)}$	<pre>Xmap=map([Np/2 2],{},0:Np-1)</pre>
Computer P _{ID} 0 1 2 3 Pid		
 A map is a mapping of array indices to processors Can be block, cyclic, block-cyclic, or block w/overlap 		

• Use P() notation (or map) to set which dimension to split among processors



A processor *map* for a numerical array is an *assignment* of *blocks* of *data* to processing elements.





pMatlab constructors are overloaded to take a map as an argument, and return a distributed array.





Advantages of Maps





Redistribution of Data



- Different distributed arrays can have different maps
- Assignment between arrays with the "=" operator causes data to be redistributed
- Underlying library determines all the message to send



- Parallel Design
- Distributed Arrays

Concurrency vs Locality

- Definition
- Example
- Metrics

- Execution
- Summary



Definitions

Parallel Concurrency

- Number of operations that can be done in parallel (i.e. no dependencies)
- Measured with:

Degrees of Parallelism

Parallel Locality

- Is the data for the operations local to the processor
- Measured with ratio: Computation/Communication
 = (Work)/(Data Moved)



- Concurrency is ubiquitous; "easy" to find
- Locality is harder to find, but is the key to performance
- Distributed arrays derive concurrency from locality



Serial

<u>Math</u>	Matlab
X,Y : $\mathbb{R}^{N \times N}$	X = zeros(N,N); Y = zeros(N,N);
for i=1:N for j=1:N Y(i,j) = X(i,j) + 1	<pre>for i=1:N for j=1:N Y(i,j) = X(i,j) + 1; end end</pre>
Concurrency: maxLocality	degrees of parallelism = N ²
 Work = N² Data Moved: dependent 	ends upon map



1D distribution

<u>Math</u>	pMatlab
$\mathbf{X}, \mathbf{Y}: \mathbb{R}^{P(N) \times N}$	<pre>XYmap = map([NP 1],{},0:Np-1); X = zeros(N,N,XYmap); Y = zeros(N,N,XYmap);</pre>
for i=1:N for j=1:N Y (i,j) = X (i,j) + 1	<pre>for i=1:N for j=1:N Y(i,j) = X(i,j) + 1; end end</pre>
 Concurrency: degrees of parallelism = min(N,N_P) 	

- Locality: Work = N², Data Moved = 0
- Computation/Communication = Work/(Data Moved) $\rightarrow \infty$



2D distribution

<u>Math</u>	<u>pMatlab</u>
X,Y : $\mathbb{R}^{P(N) \times P(N)}$	<pre>XYmap = map([Np/2 2],{},0:Np-1); X = zeros(N,N,XYmap); Y = zeros(N,N,XYmap);</pre>
for i=1:N for j=1:N Y(i,j) = X(i,j) + 1	<pre>for i=1:N for j=1:N Y(i,j) = X(i,j) + 1; end end</pre>
 Concurrency: degrees of parallelism = min(N²,N_P) 	

- Locality: Work = N², Data Moved = 0
- Computation/Communication = Work/(Data Moved) $\rightarrow \infty$



2D Explicitly Local

<u>Math</u>	<u>pMatlab</u>
X,Y : $\mathbb{R}^{P(N) \times P(N)}$	<pre>XYmap = map([Np/2 2],{},0:Np-1); Xloc = local(zeros(N,N,XYmap)); Yloc = local(zeros(N,N,XYmap));</pre>
<pre>for i=1:size(X.loc,1) for j=1:size(X.loc,2) Y.loc(i,j) = X.loc(i,j) + 1</pre>	<pre>for i=1:size(Xloc,1) for j=1:size(Xloc,2) Yloc(i,j) = Xloc(i,j) + 1; end end</pre>

- Concurrency: degrees of parallelism = min(N², N_P)
- Locality: Work = N², Data Moved = 0
- Computation/Communication = Work/(Data Moved) $\rightarrow \infty$



Math

1D with Redistribution

$ \begin{array}{l} \mathbf{X} & : \ \mathbb{R}^{\mathbf{P}(\mathbf{N})\mathbf{x}\mathbf{N}} \\ \mathbf{Y} & : \ \mathbb{R}^{\mathbf{N}\mathbf{x}\mathbf{P}(\mathbf{N})} \end{array} $	<pre>Xmap = map([Np 1],{},0:Np-1); Ymap = map([1 Np],{},0:Np-1); X = zeros(N,N,Xmap); Y = zeros(N,N,Ymap);</pre>
for i=1:N for j=1:N Y(i,j) = X(i,j) + 1	<pre>for i=1:N for j=1:N Y(i,j) = X(i,j) + 1; end end</pre>

- Concurrency: degrees of parallelism = min(N,N_P)
- Locality: Work = N², Data Moved = N²
- Computation/Communication = Work/(Data Moved) = 1



- Parallel Design
- Distributed Arrays
- Concurrency vs Locality



Summary

- Four Step Process
- Speedup
- Amdahl's Law
- Perforfmance vs Effort
- Portability



Running

- Start Matlab
 - Type: cd examples/AddOne
- Run dAddOne
 - Edit pAddOne.m and set: PARALLEL = 0;
 - Type: pRUN('pAddOne',1,{})
- **Repeat with: PARALLEL = 1**;
- Repeat with: pRUN('pAddOne',2,{});
- Repeat with: pRUN('pAddOne',2,{'cluster'});
- Four steps to taking a serial Matlab program and making it a parallel Matlab program



Simple four step process for debugging a parallel program



Always debug at earliest step possible (takes less time)



Timing

- Run dAddOne: pRUN('pAddOne',1,{'cluster'});
 - Record processing_time
- Repeat with: pRUN('pAddOne', 2, {'cluster'});
 - **Record** processing_time
- Repeat with: pRUN('pAddone', 4, {'cluster'});
 - Record processing_time
- Repeat with: pRUN('pAddone', 8, {'cluster'});
 - Record processing_time
- Repeat with: pRUN('pAddone',16,{'cluster'});
 - Record processing_time
- Run program while doubling number of processors
- Record execution time



Computing Speedup



- Speedup Formula: Speedup $(N_P) = Time(N_P=1)/Time(N_P)$
- Goal is sublinear speedup
- All programs saturate at some value of N_P



Amdahl's Law



- Serial fraction sets maximum speedup: S_{max} = w₁⁻¹
- Likewise: Speedup(N_P=w_|⁻¹) = S_{max}/2



HPC Challenge Speedup vs Effort



- Ultimate Goal is speedup with minimum effort
- HPC Challenge benchmark data shows that pMatlab can deliver high performance with a low code size



Portable Parallel Programming





- Distributed arrays eliminate most parallel coding burden
- Writing well performing programs requires expertise
- Experts rely on several key concepts
 - Concurrency vs Locality
 - Measuring Speedup
 - Amdahl's Law
- Four step process for developing programs
 - Minimizes debugging time
 - Maximizes performance



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